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Multi-Objective Artificial Bee Colony for Designing Multiple Genes Encoding the Same Protein

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Abstract

The improvement of protein expressio. I vels represents one of the most important goals in synthetic biology. A ora r to accomplish it, a promising and widely-used strategy lies on integratine multiple genes that encode the subject protein into an organism genome. L'is important task, however, is affected by several challenging issues. rin, 'ly, the integration of highly similar sequences can potentially induce ho. plogor recombination, a negative effect that implies a reduction in the nur .ber of genes effectively integrated. This is the reason why it is important to desig. mulliple protein-coding sequences (also named CDSs) that are as differ no is possible, between both different CDSs and different subsequences with in the same CDS. Additionally, codon usage frequencies in these CDSs should be is highly adapted to the organism as possible. Therefore, this task involves different and conflicting objectives that must be optimized, thus being sure ¹ e to be tackled as a multi-objective optimization problem. In this work we design and implement the algorithm MOABC (Multi-Objective Artificial Pee Colony) to solve the problem of designing multiple CDSs that encode he san protein, considering three objectives to be optimized. The experimen-

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tal evaluation herein performed suggests that MOABC is able to <u>b</u>tain elevant results, showing statistically significant improvements over the ones found in the literature.

Keywords: Multi-Objective Artificial Bee Colony, De^{c} .gn of Multiple Genes, Encoding of the Same Protein, Protein-Coding Sequence (CD^{c}).

1. Introduction

The idea of maximizing the expression levels of proteins is becoming an increasingly relevant research goal in the field of symbetic biology. In this context, the integration of multiple genes encoding a certain protein into an organism genome represents a promising approximation of high a certain protein into an organism genes encoding the same protein implies, in general terms, that the expression levels of this protein can also increase n times (see for example [1]). Although this effect does not happen in accolutery all the cases [2], such strategies have been attracting increasing interest throughout the years [3, 4, 5].

- However, this is a difficult tak. The process of integrating multiple genes into an organism genome is a polex and time-consuming, and also it involves high cost. In order a refluce all these problems, in the current protocols, the multiple genes are integrated very near each other within the organism genome [3, 4, 5]. This is a good solution, but it has a weak point. When repetitive
- sequences are very near each other, they can induce homologous recombination, and as a consequence, some of these sequences are lost [6]. For example, if a gene is replicated 4 times (g_1, g_2, g_3, g_4) and these copies are concatenated in tander and a genated into a genome, a homologous recombination could take place betwee $i g_1$ and g_4 , and the copies of gene g_2 and g_3 would be lost.
- For unat reason, it is very important that the multiple protein-coding seuences (also named CDSs) are as different as possible, between both different CDSs and different subsequences within the same CDS. The exact length that are identical sequences must have in order to induce homologous recombination is not known, and it could change depending on the organism. For example, [7]

- ²⁵ reported, in *Escherichia coli*, that the identical sequences shoun' have ut least 23 bp (base pairs). [8] reported, in *Bacillus subtilis*, homologous recombination of identical sequences with a length of 70 bp. [9] reported, 'a *Saccharomyces cerevisiae*, a highly increased rate of homologous recombination with identical sequences of 30 bp. Although the exact length is not 'rnow', all the studies
- agree that the longer the identical sequences are, the high random rando
- The way to obtain different CDSs that create the same protein is using different codons for the corresponding mino acids. Almost all the amino acids can be encoded by different synol, in our odons (a codon is a series of three nucleotides). Therefore, chancing the synonymous codons used for encoding a particular amino acid we can obtain a different CDS that encodes the same
- ⁴⁰ protein. Synonymous cod his crcur with different frequencies in an organism, and the choice of codon. may : fect the expression levels of a protein [10]. For this reason, it is important to select the codons with the highest usage frequencies (the most . only adapted ones). This will be the third objective in our multi-object. optimization problem, in which we will try to maximize
- ⁴⁵ the minimum close of the Codon Adaptation Index (CAI). As we can observe, different and cordicting objectives have to be optimized.

Regar .ing the related work, codon usage frequency optimization in an individual CL ^C nas been analyzed in recent studies [11, 12, 13], and also different tools nave been proposed, such as COOL [14], D-Tailor [15] or OPTIMIZER [16].

- In th. regar i, [17] demonstrated that the codon usage of highly expressed genes vas selected in evolution to maintain the efficiency of global protein translation,
 a. d [18] shed insight into the factors that influenced the codon usage frequency
 b. des associated with the central nervous system. However, our approach
 c. different from all these previous related works, because we also optimize se-
- $_{\tt 55}$ $\,$ quences to increase the nucleotide differences among multiple CDSs that encode

the same protein. In that sense, after a literature review, to our the standard sense, after a literature review, to our the standard sense, the only previous proposal that addresses the same multi-object re-optimization problem has been recently published [19]. In that last paper, the problem was solved by using the NSGA-II algorithm. Therefore, we compare our results with this previous related work in order to show the advantages of our approach.

In particular, our approach is based on the Multi-Collective Artificial Bee Colony (MOABC) algorithm. This algorithm is a multi-ed jective adaptation of the Artificial Bee Colony (ABC) algorithm [20], which has been selected due to its good results in other applications [21]. In the port, we have adapted the

- ABC algorithm to the multi-objective contant are we have designed and implemented the MOABC algorithm for designing a ret of CDSs that encode a protein avoiding inducing homologous recombination and using the synonymous codons with the best CAIs. Apart from the MC.'BC algorithm designed and implemented here, currently, there are a number of MOABC algorithms available in
- the literature. In this sense, [22] developed a MOABC that used a grid-based approach to adaptively assessing Pareto front maintained in an external archive. The external archive was red to control the flying behaviors of the individuals and structuring the ble colony. [23] proposed a MOABC where an elite-guided solution generation structure exploit the neighborhood of the solution generation structure.
- existing solution is sed on the guidance of the elite. Furthermore, a novel fitness calculation pethod was presented to calculate the selecting probability for onlooker bec. [2] designed a MOABC with dynamic population, which synergized the dea of extended life-cycle evolving model to balance the exploration and exploit cion trade-off. In this approach, the bee was able to reproduce and
- die dynamically throughout the foraging process and population size varied as the a. orither ran. [25] developed a MOABC that integrated genetic operators. In this way, this approach used the traditional crossover and mutation offspring process

²⁵ our best knowledge, this is the first time that a swarm intelligence al-²⁵ orithm is used to solve our multi-objective problem. As we will see, after a comparative study and the corresponding statistical evaluation, we can conclude that our approach obtains very good results. In fact, as a final ontrolation of our work, we can highlight that our approach obtains better a sults than the ones previously published in the literature.

The rest of this paper is organized as follows. Sect on 2 ϵ plains and gives a formal definition of the multi-objective optimization , robler to solve. After that, section 3 details and describes our approact (M \sim BC) to address this optimization problem. Section 4 includes the exportments performed, the results obtained, and the comparisons with the results found in the literature. Finally,

section 5 explains the conclusions of this work with indicates possible future lines.

2. Problem Definition

Given a protein to encode, a solution to vir multi-objective optimization problem is a set of sequences (CDSs, encoding this protein. The number of CDSs
per protein is determined by the user. Every CDS is a sequence of nucleotides, and therefore, we represent a solution to our problem as a set of sequences of characters. Given a pretein, all is CDSs will have the same length. An example of solution is shown in T.ole 1.

The evaluation of each solution is based on three objective functions. The first function concerns the encoding of each codon (preferring the codons with the highest v age frequencies) and the other two functions are related to the avoidance of re_{r} titions between CDSs and between subsequences within the same CUS. 7 he following subsections explain in detail these three objective functions.

110 2.1. Codor Adaptation Index (CAI)

The firs objective function is related with the occurrence frequency of each course, because some codons have higher frequency than others, and therefore, j is better to select these better adapted codons. The focus in this function

is the minimum value of Codon Adaptation Index (mCAI). Eq. tion 1 shows how this calculation is made.

$$mCAI = \min_{1 \le i \le I} \quad CAI(CDS_i) \tag{1}$$

where CDS_i is each CDS that encodes the protein I = 1 number of CDSs, and CAI value is calculated for each one as indicated by Equation 2.

$$CAI(CDS_i) = \sqrt[N]{\prod_{n=1}^{N} W(rodon_{,n})},$$
(2)

where N is the number of codons $\lim_{n \to D} O_{D}_{i}$ has and W is the weight assigned to the $codon_{i,n}$. This weight is calculated as the usage frequency of $codon_{i,n}$ relative to (divided into) the mage frequency of the most frequent codon among the synonymous codens of $codon_{i,n}$ [26]. The usage frequencies have been obtained from the resumble prried out by [19].

The objective is to optimize the minimum CAI value among all the CDSs (mCAI). The use of the alreage CAI value would not be sufficiently appropriate because it is possible that there exists a non-the CDS with a very low CAI within a good average CAI. In conclusion, the objective is to maximize the minimum CAI value (mCAI) to a bid we that all the CDSs have high codon adaptation index.

2.2. Hammi q D stance between CDSs (HD)

The second objective function is focused on selecting CDSs that are very different between the 1. This function is based on a normalized Hamming distance value between two CDSs from a solution. The objective is calculated as the minin run velue among all these pairs of CDSs (see Equation 3).

$$mHD = \min_{1 \le i < j \le I} \quad \frac{HD(CDS_i, CDS_j)}{L}.$$
(3)

For a pair of CDSs, CDS_i and CDS_j , both with length L 1. clean les, the 135 Hamming Distance (HD) is calculated as shown in Equation 4.

$$HD(CDS_i, CDS_j) = \sum_{1 \le k \le L} \sigma(CDS_{i,k}, \square DS_{j,k}),$$
(4)

where the *i*-th and *j*-th CDSs are compared σ a, in both CDSs, the *k*-th nucleotide is evaluated. If $CDS_{i,k}$ and $CDS_{j,k}$ require qual then σ will be 0. However, if they are different nucleotides, σ is set τ 1.

The objective is to maximize mHD. Again, we use the minimum value, as in the first objective function, because if we use the average value, we could have a very low HD within a good average.

2.3. Length of Repeated or Common Surveyings (LRCS)

The third objective function is focured on the idea of selecting different subsequences between different CLUS and also within the same CDS. In conclu-145 sion, this objective tries to reduce the length of repeated or common substrings

(LRCS).

We say that we find a common substring $S_{i,p,l}$ in the *i*-th CDS, at the *p*th position and with a length of *l* characters (nucleotides), when the same or another CDS (*j*-t⁺, CDS)⁻¹ as the same substring $S_{j,q,l}$ at the same (in this case, $i \neq j$) or different q-th position. For example, in Table 1, *GUGUUA* is the longest common substring between all pairs of CDSs, although there are other repeated or obstrings, e.g. in CDS_3 , UGCU, but they have a lower length. The objective is to minimize the maximum length of the repeated or common substring (μ_i, τ, I, CS) found among CDSs (see Equation 5).

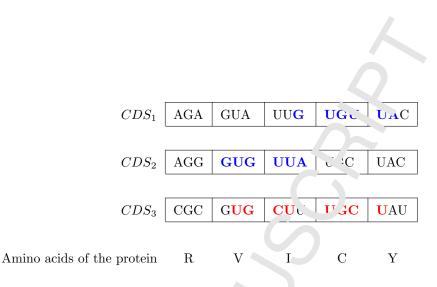


Table 1: A possible solution with 3 CDSs. Example for the computation of the length of repeated or common substrings. GUGUUA is the long. + common substring between all pairs of CDSs, although there are other repeated substrings, c.g. in CDS_3 , UGCU, but they have a lower length.

$$MLRCS = \max_{1 \le j \le j} \frac{LRCS(CDS_i, CDS_j)}{L},$$
(5)

where L is the length in nucleoticles of the CDSs. For every pair of CDSs, CDS_i and CDS_j (observe that i = j is allowed), LRCS is calculated as in Equation 6.

$$LRCS(C \supset S_i, C \supset S_j) = length(S_{i,p,l}) \quad when \ (S_{i,p,l} = S_{j,q,l}),$$
(6)

where L, as said, is be length of the CDSs, and if p = q then $i \neq j$.

3. Multi-Dbje, "ive Artificial Bee Colony (MOABC) Algorithm

Artificial L ~ Coony (ABC) is an algorithm introduced by Dervis Karaboga in 2007 [20], notivated by the intelligent behavior of honey bees and used as an optimeration tool.

ABC defines a set of operations in order to find pseudo-optimal solutions in a .imil r way to how bees find their best food sources. Individuals are members to a colony where each type of bee has a task and there are three types of bees:

• Employed bees are associated to a specific food source (solution).

- Onlooker bees watch the dances of employed bees and c. ose the most profitable food sources depending on the dances.
- Scout bees carry out random searches for discovering new food sources.

In our approach, a Multi-Objective (MO) adapta on of ABC algorithm (MOABC) has been designed and implemented to chie cood solutions for our problem. In particular, as we address a multi-conjective optimization problem, there is not one only best solution, but the result fill be a set of trade-off solutions that optimize in different ways the constant optimize.

- An important concept used in MO optimization problems is the Pareto dominance between two solutions. A solution x do. vinates another solution y (represented as $x \succ y$) when the values obtained by x in all the objective functions are always better than or equal to the corresponding values obtained by y, and at the same time, x obtains a better value in at least one of the objective functions.
- In the same way, we say that a solution is non-dominated or Pareto optimal if there is no other solution may dominates it. The graphical representation of the non-dominated solution set (or Pareto set) is known as Pareto front. The Pareto set represents the cubset of best solutions found for the multi-objective optimization problem.

Remember that we represent a solution as a set of equal-length sequences of characters (see Typle 1 for an example). Algorithm 1 shows the pseudo-code of the MOABC proposed. The first step is to establish a empty file to store non-dominate solutions. Then we initialize the colony with a total of *colony_size* solutions (". e 2) Each individual in the colony is created randomly, except one,

that is created by selecting each codon with the highest weight, and therefore, its $m \,^{\gamma} AI$ v tue will be 1. This particular solution is a non-dominated solution and co. 'd help future generations to achieve high CAI values.

Algorithm 1 MOABC pseudo-code.

Input: colony_size (number of individuals/solutions), max_cy_les (maximum number of cycles/generations), limit (abandonment cruevicu), and P_m (mutation probability)

Output: nondominated_file (file with the non-domin. *ed so' itions)

- 1: nondominated_file $\leftarrow \emptyset$
- 2: *init_colony(colony_size)*
- 3: for $cycle \leftarrow 1, max_cycles$ do
- 4: $send_employed_bees(colony_size, P_m)$
- 5: $rank_and_crowding(colony_size)$
- $6: \quad calculate_probabilities(colony_size)$
- 7: $send_onlooker_bees(colony_size, P_m)$
- 8: $send_scout_bees(limit, cycle_)$
- 9: rank_and_crowding(2 * colony__ize)
- 10: $update_nondominated_solut_{\sim} \neg s(nondominated_file)$
- 11: end for

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Next, a for loop starts, when a includes operations that make the bee colony evolve for max_cycic cycles or generations. Each search cycle involves the management of e ployed bees, onlooker bees, and scout bees.

Employed bees are the first ones to perform (line 4). Every employed bee has an associated solution and a random mutation (among the different types of mutation) is apple \neg to each individual for improving the quality. The mutation operator whenges a specific part of a solution with the goal of optimizing some of the pojective functions. The mutated solution will only be selected if it dominates the original solution. In this study, we use four types of mutation (as

sua, every time a solution is mutated, one of these mutations will be randomly . elected and applied):

. For each CDS, each codon is randomly replaced by another encoding (synonymous codon), with a probability of P_m . We can observe an example of this mutation in Figure 1.

- 2. For the CDS with the minimum CAI value, each codon \neg replaced by other encoding with greater weight (if several synon mous codons have higher weight, one of them is randomly selected), which a probability of P_m . If the codon has the encoding with the high st weight, then there is not replacement. An example of this mutation is thown in Figure 2.
- 3. For the pair of CDSs with the minimum I.D, i.e., codon is randomly replaced by another encoding, with a probability of \mathcal{P}_m . Figure 3 presents an example for this mutation.
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4. Codons that are within the longest right common substring are randomly replaced by another encodies, when a probability of P_m . Figure 4 illustrates an example of this mutation.

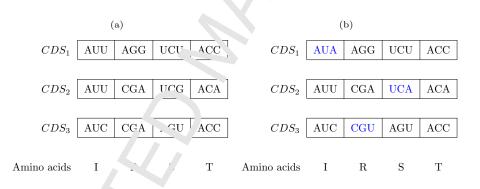


Figure 1: Example of random mutation in a solution with 3 CDSs, where (a) represents the original solution and (b) represents the mutated solution. After the mutation, the first codon in CDS_1 , \cdots third codon in CDS_2 , and the second codon in CDS_3 are replaced by synonymous codoms (in blue) randomly selected.

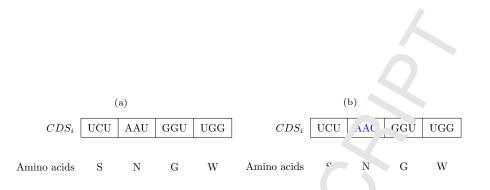


Figure 2: Example of mutation in CDS_i , which is the CF $_{\mathcal{S}}$ with the lowest CAI value in a solution with *n* CDSs. In the original solution (a), the A^{\prime} value is 0.966. Applying the mutation (b), the second codon is replaced by another ∇non sus codon (in blue) with greater weight and, after that, CDS_i has a CAI value equal to '.

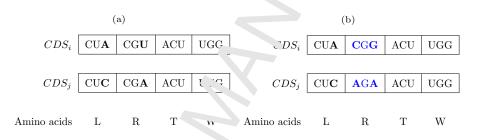


Figure 3: Example of mutation in $C_1 S_i$ and CDS_j , which are the pair of CDSs with the minimum HD in a solution w. b n CDSs. In the original CDSs, shown in (a), only two nucleotides are different (in bold) between them, so the HD is 0.167. After the mutation (b), the second codon in $-c^1$ CDS has been replaced by another synonymous codon selected at random. Consequently, the succeeding are different (in bold) and the HD has been incremented (0.25C between these CDSs.

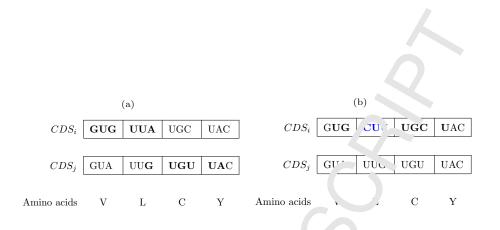


Figure 4: Example of mutation in a solution with $n \text{ CI } S_i$, where DDS_i and CDS_j are the pair of CDSs that contain the longest repeated or common substring. In the original CDSs shown in (a), the longest common substring is $GUGU \subset A$ (in ¹ old). After the mutation (b), the second codon of CDS_i has been replaced by anothe synonymous codon (in blue) selected at random. Therefore, the previous common substring has been broken and the new longest common substring is UGCU (in bold) in CDS_i , which is a shorter one.

In the onlooker bees phase (h < 7) the colony size is duplicated. Every onlooker bee chooses an employed beam as its initial solution. This selection is based on the probability value associated (line 6) with each employed bee. Employed bees that represent incluer solutions will have higher probability values. To check which employed bees are better, they are evaluated and sorted by two metrics: rank and crowding (line 5). After a non-dominated sorting, the rank value indicates in which age of the generated Pareto fronts a solution is. This non-dominated solutions. The second metric computes the solutions' crowding distance. The ground the crowding distance among the solutions, the greater

the solutions divertity. A more detailed explanation of these two metrics can be found in [27].

F ery or looker bee will try to improve the selected employed bee (its initial solut. n) by using the same random mutation procedure explained above. In nis care, the mutated solution is selected if it is non-dominated by the original core, otherwise the original solution is kept.

will not be able to be improved, it has been exhausted, so it is and and replaced with a new random solution. Furthermore, this new random solution is mutated n times, proportional to the current cycle, so that is can compete in the next cycle with the other solutions in the colony.

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Before next cycle, the colony is reduced by half (it. original size). For selecting the best solutions, these are ordered by rank and prowding again, with a difference: this process is applied to the whole colour, with double size of the original one (line 9). Finally, the non-dominated polutions in this cycle are stored in the *nondominated_file* (line 10) and the more next cycle can begin.

After explaining the MOABC pseudo code — can highlight its strengths and advantages. It is a multi-objective evolutionary algorithm based on the intelligent behavior of a swarm of bees a. donen task social cooperation. From a theoretical perspective, swarm intering once opproaches are built upon the defini-

- tion of autonomous agents with specific roles that cooperate together to achieve a common goal. These agents below of to different classes with different tasks and local rules, whose interactions lead to the attainment of global, collective intelligence through the same of information among all the components of the swarm.
- In MOABC, the division of labor is implemented through the identification of different bee classes, each one with a well-defined role. While employed bees check the neighter shood of solutions previously identified by the swarm, onlooker bees exploit the clost promising solutions found by the employed bees. On the other har 4, scout plees have the responsibility of addressing local optima issues,
- identifying tagrant solutions and performing exploration tasks to find new promising condidates in unexplored search space regions. The sharing of informatic or involves the communication of high-quality solutions from the employed bees to the onlooker ones, the identification of exhausted employed/onlooker solutions is that must be replaced by the scout bees, and the definition of the next grant ation employed solutions from all the agents in the swarm. In this way, the population is processed by the local rules of each bee class, whose interactions allow the global spread of information among all the members thus leading to

the global improvement of the population.

Please observe that, by considering different bee categories, the implemented mutation operator can be applied in different ways in accordence with the role (exploitation/exploration) of the current bee under processing. More specifically, while checking the neighborhood (employed) and exploiting high-quality solutions (onlooker), new solutions are generated by $\mathcal{P}_{a'P}$ ving once the mutation operator. On the other hand, the exploration tasks (shout) involve multiple applications of the operator over randomly generated polutions.

In comparison to other evolutionary methods, MCABC undertakes the optimization process following the task division interactions, and self-organization of the components of the bee swarm. The definition of employed, onlooker, and scout bees allows the integration of multiple search strategies (including exploitation and exploration-oriented procedures), which are applied in accordance with the current status of the production. In this sense, the inclusion of the *limit* control parameter allows the algorithm to quickly identify the presence of local optima, applying a specific technique (scout searches) to deal with them. The sharing of information betwisen bees also represents a distinctive feature

- over traditional evolution ry designs, as it allows the search engine to address the problem by considering all the information gathered by the entire swarm. All these elements is "re rise to a robust search engine that can boost the solution of complex opt" ization problems like the design of CDSs. In fact, the literature gives acrow a of the relevance of ABC in comparison to other approaches (such as genetic algorithms and differential evolution) in multiple sets of numer
 - ical test not tion [28], multi-variable [20], and multi-dimensional [29] scenarios. Improved r sults have also been reported for the case of real-world problems [21], reluding hard-to-solve problems from the biological domain [30, 31].

In this study, this pseudo-code has been implemented in C/C++. The next solution shows the data sets used in our experiments, the parameter settings for No. BC, the results obtained, the comparison with the results found in the interature, and the statistical analysis of the results.

4. Experiments and Results

for every protein.

As we have compared with the results from Terai et al. [19] in cur experiments, we used nine proteins as a representative sample based or "vo attributes: length of the protein (in AA, Amino Acids) and the number of CDSs for that protein. Since both attributes influence the complexity of the instance, we balanced these two attributes. In particular, Table 2 shows by the ender proteins have different trade-offs between length and number of CDSs, so for a larger number of CDSs, we chose a protein with a smaller length and vice versa. Observe that this table includes instances with very different number of CDSs and lengths, and also, with different complexities, being a representative set of instances. We have used the Universal Protein Reso (continent ¹) to get the FASTA format

			· · · · · · · · · · · · · · · · · · ·	
Code	Name	CDrs	Length (AA)	CDSs*Length
Q5VZP5	DUS27_HUMAN		1158	2316
A4Y1B6	FADB_SH'.PC	3	716	2148
B3LS90	OCA5_YE. S1	4	679	2716
B4TWR7	CAI7_SALSV	5	505	2525
Q91X51	GOUS1_NOUSE	6	446	2676
Q89BP2	ΓAF. BRADU	7	388	2716
A6L9J9	ı.?PF_PARD8	8	221	1768
Q88X33	▼ 1415_LACPL	9	114	1026
B7KH'.9	PETG_CYAP7	10	38	380

Table 2: List of proteins used in the experiments.

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For each instance in Table 2, we applied the algorithm explained in the vervious section and we compared with the results obtained in the method implemented by Terai et al. [19] (based on NSGA-II algorithm). For getting

¹http://www.uniprot.org/uniprot/

these results, we used the web-based application² provided by these tuthors, using their default parameters set.

- In order to provide a fair comparison, the value of the cameters for the colony size and the number of generations were set to ne same value as in the NSGA-II algorithm, that is, the colony size is 100 individuals (solutions) and the number of generations (max_cycles) is equal to 100. Ir an OABC algorithm, we adjusted other two parameters: the maximum number c, attempts (or limit)
- to improve an employed or onlooker bee (10 attem, ts in our case) and the probability of mutation (P_m) that is equal to $e^{\alpha r}$ (5%). We experimented with different values for these two parameters f find the best configuration. Specifically, Table 3 shows the values tested a. 4 the best adjustments are given by the highlighted values.

Furthermore, to ensure reliable scatters we repeated every experiment 31 times due to the stochastic nature of NOABC algorithm.

	Checked values						
P_m	1.25%	2.5%	5%	10%	20%	30%	40%
limit	ç	-	10	15			

Table 3: All tested values to indice best configuration. The best value is highlighted in bold.

To evaluate the quality of the results, we made use of two indicators widely used in multoblective optimization: hypervolume [32] and set coverage [33]. We also r alized a tatistical analysis to find the statistical significance and to be sure upt the esults are not likely to occur randomly.

4.1. Hyperve 'ume indicator

7. ne hypervolume indicator (HV), also known as Lebesgue measure [34], is a nary i dicator of the quality of a set of non-dominated solutions. In the case f 3 objective functions, this measure computes the volume (in percentage) of

²http://tandem.trahed.jp/tandem/

the objective space covered by a Pareto front \mathcal{A} with points $(a_1, a_2, \ldots, a_{|\mathcal{A}|})$, taking into account a reference point r. Equation 7 indicates how to calculate the value of HV.

$$HV(\mathcal{A}, r) = Leb\left(\bigcup_{i=1}^{|\mathcal{A}|} h(a_i, r)\right),\tag{7}$$

where *Leb* refers to the Lebesgue measure, $|\mathcal{A}|$ ne si' e (cardinality) of set \mathcal{A} , and $h(a_i, r)$ is the volume of the cube defined \mathbb{C}^* each point in \mathcal{A} (taking also into account the reference point).

Table 4 indicates the nadir and ideal values used in the hypervolume computations for all the proteins. These values were obtained by considering the results in all our experiments. Taking use a count these values, the hypervolume calculations have been performed on probjective scores normalized in the scale [0,1] to avoid the influence of dufferent ranges in the objective values.

Objective	Nadı. ² value	Ideal value	
mCA 1	0	1	
mHD	0	0.35	
N.LR JS	1	0	

Table 4: Nadir an , idea. ralues used in the hypervolume computations and normalizations for all the protei ...

The H \checkmark indicator has been calculated by the same way for both algorithms. Table 5 shows modian HV results and their quartile deviations calculated after assessing each protein individually. In almost all the proteins, the HV values obtailed by the MOABC algorithm are better than those from the method proposed low Terai et al. [19], with the exception of the protein Q5VZP5 (as we will see in this case, the differences between both algorithms are not statisticomposition. Finally, the last row shows the average value among the nine proteins. Again, we can conclude that the MOABC algorithm is better than the NSGA-II algorithm [19]. This means that MOABC is able to obtain better

Protein	MOABC	NSGA II [19]
Q5VZP5	$68.43\%_{\pm 0.62\%}$	68.48 $\%_{\pm 0.002}$ %
A4Y1B6	$60.35\%_{\pm 0.26\%}$	6° JJ %+0.0007%
B3LS90	$63.73\%_{\pm 0.14\%}$	6. 1. M±r 0018%
B4TWR7	${f 57.06\%_{\pm 0.15\%}}$	$55.^{9}\%_{\pm 0.0024\%}$
Q91X51	$59.71\%_{\pm 0.19\%}$	$57.6^{5} \ \%_{\pm 0.0026\%}$
Q89BP2	$57.42\%_{\pm 0.15\%}$	$55.55\%_{\pm 0.0025\%}$
A6L9J9	$53.77\%_{\pm 0.16\%}$	$52.07\%_{\pm 0.0015\%}$
Q88X33	$48.71\%_{\pm 0.2}$ %	$46.93\%_{\pm 0.0010\%}$
B7KHU9	47.71% 54%	$43.72\%_{\pm 0.0014\%}$
Average	57.43%	55.91%

Pareto fronts, which include better non-dominated solutions that "over thigher volume of the objective space.

Table 5: Results for hypervolur $\sim \dots$ 'cator, in the format: median $\pm quartile_deviation$. In bold we highlight the better result \sim

In addition, Figure 5 shows a visual representation for one of the proteins, B7KHU9. The objective values are already normalized in the scale [0,1] taking into account Table 4. As can be seen from that graph, the points in the projection (MLRCS, mHD) are near for both algorithms. By contrast, for the projection (MLRCS, mHD) and (mCAI, MLRCS), in most of the points, the solution from MCABC are better than the solutions from NSGA-II. Moreover, MOAPC cover, more regions of the objective space than NSGA-II. These improvements explain the clear advantage of MOABC with respect to NSGA-II.

 F_{24} F is the rest of the analyzed proteins, the behavior is similar.

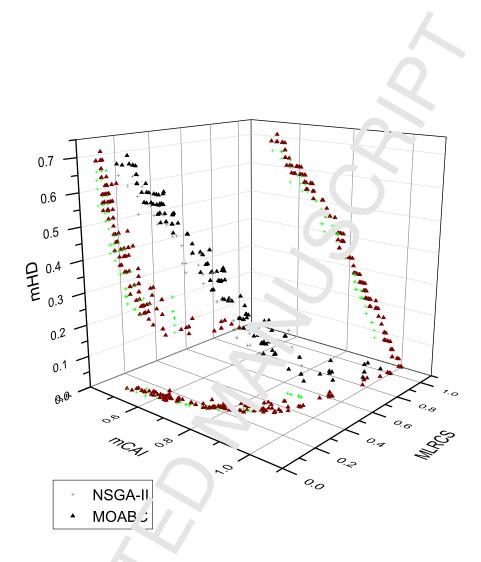


Figure 5: 3D scatt r pt. of the median Pareto fronts for the B7KHU9 protein. The points in the different 2^{r} , rojections appear in red (MOABC) or green (NSGA-II), using the corresponding symbol.

4.2. Set Con rage indicator

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Set c verage $(\Box C)$ is the second quality indicator used. In this case, it is a bina. v-type ndicator. This measure is based on how many solutions/points l elonging to a Pareto front \mathcal{B} are covered by solutions/points from another Pareto f ont \mathcal{A} . As Equation 8 shows, a solution b_j is said to be covered if there is a solution a_i that dominates to b_j or is equal to this one.

$$SC(\mathcal{A}, \mathcal{B}) = \frac{|\{b_j \in \mathcal{B}; \exists a_i \in \mathcal{A} : a_i \succeq b_j\}|}{|\mathcal{B}|},\tag{8}$$

where $|\mathcal{B}|$ is the size (cardinality) of set \mathcal{B} .

If each solution in \mathcal{B} is covered at least by a solution in \mathcal{A} then $SC(\mathcal{A},\mathcal{B})$ is equal to 1. Otherwise, if none of the solutions belonging to \mathcal{B} , solvered by the so-

- ³⁷⁵ lutions in \mathcal{A} , SC(\mathcal{A} , \mathcal{B}) will be 0. This weak-dominance corrator is not symmetric and it can happen that $SC(\mathcal{B}, \mathcal{A}) \neq 1 - SC(\mathcal{A}, \mathcal{B})$. For this reason, we calculated the measure SC in both directions for each protein SCC. (NSGA-II) and SC(NSGA-II,MOABC).
- In particular, Table 6 shows the results obtained and we can say that in almost all the cases, the MOABC algorithm "tan," proved set coverage scores than NSGA-II, except in the protein B^{ATWD7} , here they have close values. This means that, for each instance, there are many more points of the Pareto front from MOABC algorithm that do not ate points of the Pareto front from NSGA-II algorithm than vice versa. A main, in average, the solutions of MOABC
- algorithm are clearly better, covering on important percentage of the solutions obtained by the method proposed by Terai et al. [19].

Protein	SC(MOABC,	SC(NSGA-II
Protein	A-II [19])، المركزية	[19], MOABC)
Q5VZP.	28.00%	14.87%
A4Y 86	54.00%	4.26%
B3LS90	34.00%	14.62%
7.4TV/R7	23.00%	$\mathbf{24.86\%}$
Q91.751	$\boldsymbol{28.00\%}$	6.63%
Q89',P2	53.00%	0.64%
AJL9J9	36.00%	5.43%
	54.00%	0.62%
B7KHU9	62.00%	7.75%
Average	41.33%	8.85%

Table 6: Results for Set Coverage indicator. In bold we highlight the better results.

4.3. Statistical significance

In order to detect if there is a statistical significance in the esume betained, we performed a statistical analysis using a significance level $(p-v, \neg e)$ of 0.05 (5%)

or confidence level of 95%. A detailed explanation of al' the sta 'stical tests used can be found in [35]. We try to apply a parametric and 'usis such as ANalysis Of VAriance (ANOVA) but before we should ensite the samples follow a normal distribution and they have homogeneous variances. For this statistical study we use the results from the hypervolume indica or.

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Firstly, the null hypothesis to be tested by the Kolmogorov-Smirnov test (KS-test) is that the sample has a norm insurroution. As we can see in Table 7, none of the cases rejects this null hypothesis, therefore, we can say that all samples have a normal distribution wit. • confidence level of 95%.

Protein	K.	р 9	
Protein	МОАЬ ~	NSGAII [19]	Pass?
Q5VZP5	0.122	0.200	Yes
A4Y1B6	0.200	0.200	Yes
B3LS90	0.0.3	0.147	Yes
B4TWR7	¢.200	0.200	Yes
Q91X'	0.200	0.200	Yes
Q89PP2	0.200	0.200	Yes
<i>F</i> ;L97 э	0.200	0.200	Yes
Q88X5.	0.187	0.200	Yes
Р, KH J9	0.180	0.200	Yes

Table 7: Normality analysis using Kolmogorov-Smirnov test.

The following test is Levene test to check if the two samples (from both providence) have homogeneous variances (null hypothesis). The results of this est for the 9 proteins are shown in Table 8. Three cases rejected this null hypothesis so this means that these ones cannot be tested by ANOVA and are

Protein	Levene test	ANOVA	Mann- Whitne,	S atistical significance
Q5VZP5	0.149	0.907	U t .ov	No
A4Y1B6	0.020	_	0.1.1	No
B3LS90	0.457	0.000	-	Yes
B4TWR7	0.313	0.000	-	Yes
Q91X51	0.147	0.000	-	Yes
Q89BP2	0.019	-	0.000	Yes
A6L9J9	0.400	000	_	Yes
Q88X33	0.068	0.6 20	_	Yes
B7KHU9	0.015	-	0.000	Yes

analyzed by a nonparametric test such as the Mann-Whitney U ost. All other cases are tested by ANOVA.

Table 8: Results of the Lev γe test to check the homogeneous variances and results of the ANOVA or the Mann-W) they $\psi \gamma e$ it depending on the previous tests to find statistical significance.

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Thus, the null hypothesis to be tested by ANOVA and the Mann-Whitney U test is that oo h samples of the same experiment are similar and there are no statistican, cognificant differences between them. As Table 8 shows, seven of nine er periments reject this null hypothesis, therefore, we can say that these seven experiments have a statistical significance with a confidence interval of 95%. By contrast two of the experiments, Q5VZP5 and A4Y1B6, accept the null hypothesis and this implies that there are no statistically significant differences by tween the tested samples.

^{A11} in all, these results match the hypervolume values shown in Table 5. In
^C creal terms, our method accomplishes a more satisfying behavior than NSGA⁴¹⁵ 1. attaining statistically significant improvements in most of the evaluation sce-

narios. Particularly, the MOABC algorithm achieves Pareto free +s (secutions) different and better than the NSGA-II algorithm [19] in 7 c at c° 9 instances.

5. Conclusion and Future Work

In this work, we have proposed a novel method for '... multi-objective design of 420 CDSs encoding the same protein, which is an important t sk in bioinformatics (and more specifically, in synthetic biology). The tockied multiobjective problem involves three fundamental objective functions ('.'AI, HD, LRCS). In our proposal, we have adapted the ABC algorithmetic to the multi-objective context and we have designed and implemented the MOABC (Multi-Objective Artificial

- Bee Colony) algorithm for doing this when conclusion, we propose to use the MOABC algorithm for achieving a range of solutions that best encode a protein with several CDSs, taking into account that the nucleotide sequences should be as different as possible (between a protein unferent CDSs and different subsequences within the same CDS, thus avoiding the homologous recombination) and at the
- same time the codon ade ptation indexes should be as high as possible. To get solutions that allow up to n. ¹/r fair comparisons with other techniques from the literature (NSG.'-II) we used the same codon usage frequencies, and both methods used the same contrary/population size and number of generations. The experiments have been "one over 9 real protein instances. These instances com-
- bine different leng hs and number of CDSs, therefore, being a representative set of instances. '1... results show that MOABC achieves better Pareto fronts (solutions) nar the results previously published in the literature in almost all the instances, being not statistically significant the differences in the only instance where MOA C obtained little bit worse results.
- As have results work, we have planned to use other alternative multi-objective deport has to compare them with MOABC, allowing us to further evaluate the good quality of the MOABC results or even improving these results. Moreover, the number of existing MOABC algorithms is high (e.g. [22, 23, 24, 25]), therefore, the comparison of several of them in the problem under study is of interest

for a next research work. This will imply their design for this $s_{\rm F}$ -rific $_{\rm F}$ oblem, their implementation, their execution, and finally, their composition. On the other hand, due to the good results obtained by MOABC, ~ intend to apply this multi-objective algorithm to other bioinformatics r ulti-of ective problems, assessing if MOABC also obtains good results in these ther r oblems.

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Highlights

- The design of multiple genes encoding the same protein is an important task

- This task can be tackled as a multi-objective opt mization problem with 3 objectives

- We have designed and implemented a solution procedure based on the MOABC algorithm

- The experiments have been done over 9 . al protein instances

- MOABC obtains better results than the ones found in the literature

